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# Connectivity of $k$ -extendable graphs with large $k$

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## Abstract

Let  $G$  be a simple connected graph on  $2n$  vertices with perfect matching. For a given positive integer  $k$  ( $0 \leq k \leq n - 1$ ),  $G$  is  $k$ -extendable if any matching of size  $k$  in  $G$  is contained in a perfect matching of  $G$ . It is proved that if  $G$  is a  $k$ -extendable graph on  $2n$  vertices with  $k \geq n/2$ , then either  $G$  is bipartite or the connectivity of  $G$  is at least  $2k$ . As a corollary, we show that if  $G$  is a maximal  $k$ -extendable graph on  $2n$  vertices with  $n + 2 \leq 2k + 1$ , then  $G$  is  $K_{n,n}$  if  $k + 1 \leq \delta \leq n$  and  $G$  is  $K_{2n}$  if  $2k + 1 \leq \delta \leq 2n - 1$ . Moreover, if  $G$  is a minimal  $k$ -extendable graph on  $2n$  vertices with  $n + 1 \leq 2k + 1$  and  $k + 1 \leq \delta \leq n$  then the minimum degree of  $G$  is  $k + 1$ . We also discuss the relationship between the  $k$ -extendable graphs and the Hamiltonian graphs.

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## 1. Introduction and terminology

All graphs considered in this paper are finite, undirected and simple. For the terminology and notation not defined in this paper, the reader is referred to [4].

Let  $G$  and  $H$  be two graphs. Let  $kH$  denote  $k$  disjoint copies of  $H$  and  $G + H$  denote the union of  $G$  and  $H$  with each vertex of  $G$  joining to every vertex of  $H$ .

A graph  $G$  is said to be *factor-critical* if  $G - v$  has a perfect matching for each  $v \in V(G)$ . Let  $G$  be a graph with a perfect matching. Then  $G$  is said to be  $k$ -extendable for  $0 \leq k \leq (v - 2)/2$  if any matching in  $G$  of size  $k$  is contained in a perfect matching of  $G$ . And  $G$  is said to be *maximal  $k$ -extendable* if  $G$  is  $k$ -extendable and for each

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$e \in E(\bar{G})$ , where  $\bar{G}$  is the complement of  $G$ ,  $G \cup \{e\}$  is not  $k$ -extendable. And  $G$  is said to be *minimal  $k$ -extendable* if  $G$  is  $k$ -extendable and for each  $e \in E(G)$ ,  $G - e$  is not  $k$ -extendable.

The concept of  $k$ -extendable graphs was introduced by Plummer [7] in 1980. Since then, extensive researches on this topic have been done (see [1,2,6–10]). In [2], Ananchuen and Caccetta proved the following result about the minimum degree of  $k$ -extendable graphs.

**Lemma 1** (Ananchuen and Caccetta [2]). *Suppose  $1 \leq k \leq (v-2)/2$  and  $|V(G)| = v$ . Then if  $G$  is  $k$ -extendable, then either  $k+1 \leq \delta \leq v/2$  or  $2k+1 \leq \delta \leq v-1$ .*

For each value of  $\delta$  given in Lemma 1, there exist  $k$ -extendable graphs with the minimum degree  $\delta$ . However, the problem that which value in these ranges is attainable for maximal  $k$ -extendable graphs remains open. Plummer [9] proposed the following problem.

**Problem 1.** *Suppose  $1 \leq k \leq (v-2)/2$  and  $k+1 \leq j \leq v/2$  or  $2k+1 \leq j \leq v-1$ . Then which  $k$ -extendable graphs having minimum degree  $j$  are maximal  $k$ -extendable?*

Motivated by this problem, we study the  $k$ -extendable graphs with  $k \geq v/4$ , that is  $v/2 + 1 \leq 2k + 1$ , which means the two intervals for  $\delta$  in Lemma 1 are separated. We prove that if  $G$  is a  $k$ -extendable graph with  $k \geq v/4$ , then either  $G$  is bipartite or  $\kappa(G) \geq 2k$ . As corollaries, we characterize the maximal  $k$ -extendable graphs with  $v/2 + 2 \leq 2k + 1$  and we show that the minimum degree of a minimal  $k$ -extendable graph with  $v/2 + 1 \leq 2k + 1$  and with  $k + 1 \leq \delta \leq v/2$  is  $k + 1$ . Also we prove that a  $k$ -extendable graph with  $k \geq v/4$  is Hamiltonian, which shows the relation between  $k$ -extendable graphs and Hamiltonian graphs.

## 2. Main result

We start this section with a few basic lemmas on  $k$ -extendable graphs.

**Lemma 2** (Yu [10]). *A graph  $G$  is  $k$ -extendable if and only if for any matching  $M$  of size  $r$  in  $G$  ( $1 \leq r \leq k$ ),  $G - V(M)$  is  $(k-r)$ -extendable.*

**Lemma 3** (Yu [10]). *Let  $G$  be a connected  $k$ -extendable non-bipartite graph. Then for each edge  $e \in E(\bar{G})$ ,  $G + e$  is  $(k-1)$ -extendable.*

**Lemma 4** (Plummer [7]). *If  $G$  is  $k$ -extendable, then  $\kappa(G) \geq k + 1$ .*

**Lemma 5.** *Let  $G$  be a graph and  $S \subseteq V(G)$ . If the size of a maximum matching of  $G - S$  is  $m$ , then the size of a maximum matching of  $G$  is at most  $m + |S|$ .*

**Proof.** Obvious.  $\square$

We need the following lemma to prove our main result, this lemma itself may serve as a useful tool in other research on matching theory.

**Lemma 6.** *Let  $G$  be a graph with order  $v = 2r + m$ . If  $G$  has a matching of size  $r$  and deleting any vertex from  $G$ , the resulting graph still has a matching of size  $r$ , then  $G$  has a matching of size  $r + 1$  unless  $G$  has exactly  $m$  odd components and no even components and each odd component is factor-critical.*

**Proof.** Suppose that the maximum matchings of  $G$  have size  $r$ . Then by Berge's formula on maximum matching, there exists a set  $S \subseteq V(G)$  such that  $o(G-S) - |S| = m$ . If  $S \neq \emptyset$ , let  $v \in S$ ,  $G' = G - v$  and  $S' = S \setminus \{v\}$ . Then  $o(G' - S') - |S'| = o(G - S) - |S| + 1 = m + 1$ . So the maximum matching in  $G'$  has size at most  $(|V(G')| - (o(G' - S') - |S'|))/2 = (2r + m - 1 - (m + 1))/2 = r - 1$ , contradicting to the hypothesis that deleting any vertex from  $G$  the resulting graph still has a matching of size  $r$ . So  $S = \emptyset$  and  $G$  has exactly  $m$  odd components. If  $G$  has an even component  $C$ , deleting a vertex  $v$  from  $C$ ,  $G - v$  has a maximum matching of size less than  $r$  since there is a vertex in each of the  $m + 1$  odd components which is not covered by the maximum matching and also  $v$  is not covered by the maximum matching. Hence,  $G$  has no even component. But deleting any vertex  $v$  from each odd component  $C$  of  $G$ ,  $C - v$  must have a perfect matching, otherwise by counting the number of vertices of  $G$ ,  $G - v$  has no matching of size  $r$ . So each component of  $G$  is factor-critical.  $\square$

Now we give the proof of our main result.

**Theorem 7.** *If  $G$  is a  $k$ -extendable graph on  $v$  vertices with  $k \geq v/4$ , then either  $G$  is bipartite or  $\kappa(G) \geq 2k$ .*

**Proof.** By contradiction. Suppose that  $G$  is a connected  $k$ -extendable graph with connectivity at most  $2k - 1$  but not bipartite. Let  $S$  be a minimum cutset of  $G$  and let  $M$  be a maximum matching in  $G[S]$ . Let  $T = S \setminus V(M)$  and  $r = |M|$ . Since  $|S| \leq 2k - 1$ ,  $|M| \leq k - 1$ . By Lemmas 2 and 4,  $G - V(M)$  is  $(k - r + 1)$ -connected. Then we have

$$|T| \geq k - r + 1 \geq 2 \quad (1)$$

and we have  $2k - 1 \geq 2r + |T| \geq k + r + 1$ , so

$$r \leq k - 2. \quad (2)$$

**Claim 1.** *For every perfect matching  $F$  containing  $M$ ,  $F \cap E(G - S)$  is a maximum matching in  $G - S$  and  $|F \cap E(G - S)| \leq k - 1$ .*

Since  $T$  is an independent set of  $G$ , by (1) and assumption that  $|V(G)| \leq 4k$ ,

$$\begin{aligned} |F \cap E(G - S)| &= (|V(G)| - 2|M| - 2|T|)/2 \\ &= |V(G)|/2 - r - |T| \leq 2k - (k + 1) = k - 1. \end{aligned}$$

If  $F \cap E(G - S)$  is not a maximum matching in  $G - S$ , then there is a matching  $F_1$  in  $G - S$  such that  $|F_1| = |F \cap E(G - S)| + 1 \leq k$ . But by Lemma 5, the size of a maximum matching in  $G - V(F_1)$  is at most

$$|V(G - S - V(F_1))| + |M| \leq |V(G)|/2 - |F_1| - 1,$$

hence  $G - V(F_1)$  does not have perfect matching, this contradicts the  $k$ -extendability of  $G$ . The proof of Claim 1 is complete.  $\square$

By Claim 1 and the fact that  $T$  is an independent set of  $G$ , we easily prove the following claim.

**Claim 2.** *The size of every maximum matching in  $G - S$  is  $|V(G)|/2 - |M| - |T|$ .*

By (1), there are two distinct vertices  $x$  and  $y$  in  $T$ . By Lemma 3, the graph  $H = G + xy$  is  $(k - 1)$ -extendable. By (2),  $M_1 = M \cup \{xy\}$  is a matching in  $H$  which has size at most  $k - 1$ . Then  $H - V(M_1)$  has a perfect matching  $M^*$  and  $M^*$  matches each vertex of  $T \setminus \{x, y\}$  to a vertex in  $V(G - S)$ . Hence,  $M^* \cap E(G - S)$  is a matching of size  $|V(G)|/2 - |M| - |T| + 1$  in  $G - S$ . This contradicts Claim 2. The proof of Theorem 7 is complete.  $\square$

**Remark 1.** The lower bound on connectivity in Theorem 7 is best possible. Let  $H_1 = K_{2k}$ ,  $H_2 = K_r$  and  $H_3 = K_s$  with  $4 \leq r + s \leq 2k - 2$  and both  $r$  and  $s$  being positive even integers. Then  $G = H_1 + (H_2 \cup H_3)$  is  $k$ -extendable but with  $\kappa(G) = 2k$ . Also the lower bound on  $k$  in Theorem 7 is best possible. The hypothesis  $k \geq v/4$  is equivalent to  $v \leq 4k$ . Let  $H_1 = \bar{K}_{k+1}$ ,  $H_2 = K_{k+1}$  and  $H_3 = K_{2k}$ , where  $\bar{K}_{k+1}$  is the complement of  $K_{k+1}$ . Then  $G = H_1 + (H_2 \cup H_3)$  is a  $k$ -extendable graph with  $v = 4k + 2$  that is not bipartite but has connectivity  $k + 1$ .

### 3. Maximal $k$ -extendable graphs with large $k$

In this section, we characterize all maximal  $k$ -extendable graphs with  $v/2 + 2 \leq 2k + 1$ . Then we show some maximal  $k$ -extendable graphs with  $2k + 1 \leq v/2 + 1$  and with  $\delta \geq v/2$ . Our results partially answer Problem 1.

**Lemma 8** (Ananchuen and Caccetta [1]). *If  $G \neq K_v$  is a maximal  $k$ -extendable graph on  $v$  vertices, then*

- (a) *if  $v/2 < 2k$ , then  $\delta \leq v/2$ , while*
- (b) *if  $v/2 \geq 2k$ , then  $\delta \leq v/2 + 2\lfloor (k - 1)/2 \rfloor$ .*

**Lemma 9** (Plummer [8] and Yu [10]). *If  $G = (X, Y) \neq K_{n,n}$  is a connected  $k$ -extendable bipartite graph and  $e = xy \in E(\bar{G})$ , where  $x \in X$  and  $y \in Y$ , then  $G \cup \{e\}$  is also  $k$ -extendable.*

**Corollary 10.** *Let  $G$  be a maximal  $k$ -extendable graph on  $v$  vertices with  $v/2 + 2 \leq 2k + 1$ . Then*

- (a) *if  $k + 1 \leq \delta \leq v/2$ , then  $G$  is  $K_{v/2, v/2}$  and hence  $\delta = v/2$ ;*
- (b) *if  $2k + 1 \leq \delta \leq v - 1$ , then  $G$  is  $K_v$  and hence  $\delta = v - 1$ .*

**Proof.** By Theorem 7, if  $k + 1 \leq \delta \leq v/2$ , then  $G$  is bipartite. Otherwise  $\delta(G) \geq \kappa(G) \geq 2k$ . When  $v/2 + 2 \leq 2k + 1$ ,  $\delta(G) \neq 2k$  by Lemma 1. Hence,  $\delta(G) \geq 2k + 1 \geq v/2 + 2$  and  $G$  is non-bipartite. By Lemma 9, we have conclusion (a). By Lemma 8(a), we have conclusion (b).  $\square$

**Remark 2.** Corollary 10 characterizes all maximal  $k$ -extendable graphs with  $v < 4k$ . It shows that the minimum degree of a maximal  $k$ -extendable graph  $G$  with  $v \leq 4k - 2$  is either  $v/2$  or  $v - 1$ . But for the case of  $v \geq 4k$ , we give a family of maximal  $k$ -extendable graphs to show that the minimum degree of  $G$  can be much more diverse.

Let  $G_i = K_{r_i}$ ,  $i = 1, 2, \dots, m$ , where each  $r_i$  is an odd number and  $r_1 + r_2 + \dots + r_m = 2k - 2 + m$ . Let  $H_j = K_{s_j}$ ,  $j = 1, 2, \dots, m$ , where each  $s_j$  is an odd number and  $s_1 + s_2 + \dots + s_m = 2k - 2 + m$ . And let  $G = (G_1 \cup G_2 \cup \dots \cup G_m) + (H_1 \cup H_2 \cup \dots \cup H_m)$ . Then it is not too difficult to verify that  $G$  is maximal  $k$ -extendable but not  $(k + 1)$ -extendable. When we take  $m = 2$ , by choosing proper  $r_i$  and  $s_i$  ( $i = 1, 2$ ), we have  $\delta(G) = t$  for all even numbers  $t$  such that  $v/2 \leq t \leq v/2 + 2 \lfloor (k - 1)/2 \rfloor$ . When we take  $m = 3$ , by choosing proper  $r_i$  and  $s_i$  ( $i = 1, 2, 3$ ), we have  $\delta(G) = t$  for all odd numbers  $t$  such that  $v/2 \leq t \leq v/2 + \lfloor (2k + 1)/3 \rfloor - 1$ .

#### 4. Minimal $k$ -extendable graphs with large $k$

In this section, we show that the minimum degree of a minimal  $k$ -extendable graph with  $v \leq 4k$  and  $k + 1 \leq \delta \leq v/2$  is  $k + 1$ . We introduce a result of Lou [6] as a lemma.

**Lemma 11** (Lou [6]). *If  $G$  is a minimal  $k$ -extendable bipartite graph, then  $\delta(G) = k + 1$ , and furthermore, there are at least  $2k + 2$  vertices of degree  $k + 1$  in  $G$ .*

**Corollary 12.** *Let  $G$  be a minimal  $k$ -extendable graph on  $v$  vertices with  $v/2 + 1 \leq 2k + 1$ . If  $k + 1 \leq \delta(G) \leq v/2$ , then  $\delta(G) = k + 1$ . Furthermore, there are at least  $2k + 2$  vertices of degree  $k + 1$  in  $G$ .*

**Proof.** By Theorem 7, if  $k + 1 \leq \delta(G) \leq v/2$ , then  $G$  is bipartite. By Lemma 11, the result follows.  $\square$

Since a  $k$ -extendable graph with  $k \geq v/4$  is rather dense, we make the following conjectures.

**Conjecture 1.** *Let  $G$  be a minimal  $k$ -extendable graph on  $v$  vertices with  $v/2 + 1 \leq 2k + 1$ . Then  $\delta(G) = k + 1$ ,  $2k$  or  $2k + 1$ .*

In particular, for the case of  $v \leq 4k - 2$ , we have the following conjecture.

**Conjecture 2.** *Let  $G$  be a minimal  $k$ -extendable graph on  $v$  vertices with  $v/2 + 2 \leq 2k + 1$ . If  $2k + 1 \leq \delta \leq v - 1$ , then  $\delta(G) = 2k + 1$ .*

## 5. Hamiltonicity of $k$ -extendable graphs with large $k$

In this section, we show that a  $k$ -extendable graph is Hamiltonian if  $k$  is sufficiently large with respect to its order.

**Lemma 13** (Dirac [5]). *If  $\delta(G) \geq v/2$ , then  $G$  is Hamiltonian.*

**Lemma 14** (Jackson [3]). *Let  $G = (X, Y)$  be a connected bipartite graph with  $|X| = |Y| = n$ . If  $\delta(G) \geq (n + 1)/2$ , then  $G$  is Hamiltonian.*

**Corollary 15.** *If  $G$  is a  $k$ -extendable graph with  $k \geq v/4$ , then  $G$  is Hamiltonian.*

**Proof.** By Theorem 7, if  $k + 1 \leq \delta(G) \leq v/2$ ,  $G = (X, Y)$  is bipartite with  $|X| = |Y| = v/2 \leq 2k$ . However,  $\delta(G) \geq k + 1 = (2k + 2)/2 > (|X| + 1)/2$ , by Lemma 14,  $G$  is Hamiltonian. Otherwise  $\delta(G) \geq \kappa(G) \geq 2k \geq v/2$ , by Lemma 13,  $G$  is Hamiltonian.  $\square$

**Remark 3.** Although we did not find new Hamiltonian graphs in Corollary 15, we did show the relation between  $k$ -extendable graphs and Hamiltonian graphs that a  $k$ -extendable graph with sufficiently large  $k$  with respect to the order  $v(G)$  is Hamiltonian. In fact, we suspect that the lower bound on  $k$  in Corollary 15 is not best possible. And hence, we give the following conjecture.

**Conjecture 3.** *If  $G$  is a  $k$ -extendable graph with  $k > (v - 2)/6$ , then  $G$  is Hamiltonian.*

The lower bound on  $k$  in Conjecture 3 is best possible. Let  $S = \{v_1, v_2, \dots, v_{2k}\}$  be an independent set and  $H = (2k + 1)K_2$  with  $V(H) \cap S = \emptyset$ . Then  $G = S + H$  is a  $k$ -extendable graph but  $G$  is not Hamiltonian as  $G$  is not 1-tough. Here  $v(G) = 6k + 2$ , that is  $k = (v - 2)/6$ . The above counterexamples also show that a  $k$ -extendable graph with arbitrarily large  $k$  (but  $v$  is also sufficiently large) is not guaranteed to be Hamiltonian.

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