

# The $k$ -proper index of graphs\*

Lin Chen, Xueliang Li, Jinfeng Liu

Center for Combinatorics and LPMC  
Nankai University, Tianjin 300071, P.R. China

E-mail: chenlin1120120012@126.com, lxl@nankai.edu.cn, ljinfeng709@163.com

## Abstract

A tree  $T$  in an edge-colored graph is a *proper tree* if any two adjacent edges of  $T$  are colored with different colors. Let  $G$  be a graph of order  $n$  and  $k$  be a fixed integer with  $2 \leq k \leq n$ . For a vertex set  $S \subseteq V(G)$ , a tree containing the vertices of  $S$  in  $G$  is called an  *$S$ -tree*. An edge-coloring of  $G$  is called a  *$k$ -proper coloring* if for every set  $S$  of  $k$  vertices in  $G$ , there exists a proper  $S$ -tree in  $G$ . The  *$k$ -proper index* of a nontrivial connected graph  $G$ , denoted by  $px_k(G)$ , is the smallest number of colors needed in a  $k$ -proper coloring of  $G$ . In this paper, we state some simple observations about  $px_k(G)$  for a nontrivial connected graph  $G$ . Meanwhile, the  $k$ -proper indices of some special graphs are determined, and for every pair of positive integers  $a, b$  with  $2 \leq a \leq b$ , a connected graph  $G$  with  $px_k(G) = a$  and  $rx_k(G) = b$  is constructed for each integer  $k$  with  $3 \leq k \leq n$ . Also, we characterize the graphs with  $k$ -proper index  $n - 1$  and  $n - 2$ , respectively.

**Keywords:** coloring of graphs,  $k$ -proper index, characterization of graphs

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## 1. Introduction

In this paper, all graphs under our consideration are finite, undirected, connected and simple. For more notation and terminology that will be used in the sequel, we refer to [2], unless otherwise stated.

In 2008, Chartrand et al. [8] first introduced the concept of rainbow connection. Let  $G$  be a nontrivial connected graph on which an edge-coloring  $c : E(G) \rightarrow \{1, 2, \dots, k\} (k \in \mathbb{N})$  is defined, where adjacent edges may be colored with the same color. For any two vertices  $u$  and  $v$  of  $G$ , a path in  $G$

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connecting  $u$  and  $v$  is abbreviated as a  $uv$ -path. A  $uv$ -path  $P$  is a *rainbow  $uv$ -path* if no two edges of  $P$  are colored with the same color. The graph  $G$  is *rainbow connected* (with respect to  $c$ ) if  $G$  contains a rainbow  $uv$ -path for every two vertices  $u$  and  $v$ , and the coloring  $c$  is called a *rainbow coloring* of  $G$ . If  $k$  colors are used, then  $c$  is a *rainbow  $k$ -coloring*. The minimum  $k$  for which there exists a rainbow  $k$ -coloring of the edges of  $G$  is the *rainbow connection number* of  $G$ , denoted by  $rc(G)$ . The topic of rainbow connection is especially meaningful and numerous relevant papers have been written. For more details see a survey [23] and a book [24].

Subsequently, a series of generalizations of rainbow connection number were proposed. The  $k$ -rainbow index is one of them. An edge-colored tree  $T$  is a *rainbow tree* if no two edges of  $T$  are assigned the same color. Let  $G$  be a nontrivial connected graph of order  $n$  and let  $k$  be an integer with  $2 \leq k \leq n$ . A  *$k$ -rainbow coloring* of  $G$  is an edge coloring of  $G$  having the property that for every set  $S$  of  $k$  vertices of  $G$ , there exists a rainbow tree  $T$  in  $G$  such that  $S \subseteq V(T)$ . The minimum number of colors needed in a  $k$ -rainbow coloring of  $G$  is the  *$k$ -rainbow index* of  $G$ . These concepts were introduced by Chartrand et al. in [9], and were further studied in [4, 5, 10, 21, 22, 26].

In addition, a natural extension of the rainbow connection number is the proper connection number, which was introduced by Borozan et al. in [3]. A path in an edge-colored graph is said to be *properly edge-colored* (or *proper*), if every two adjacent edges on the path differ in color. An edge-colored graph  $G$  is  *$k$ -proper connected* if any two vertices are connected by  $k$  internally pairwise vertex-disjoint proper paths. The  *$k$ -proper connection number* of a  $k$ -connected graph  $G$ , denoted by  $pc_k(G)$ , is defined as the smallest number of colors that are needed in order to make  $G$   $k$ -proper connected. In particular, when  $k = 1$ , the 1-proper connection number is abbreviated as proper connection number and written as  $pc(G)$ . For more results, we refer to [1, 12, 13, 14, 15, 18, 25].

Inspired by the  $k$ -rainbow index and the proper connection number, a natural idea is to introduce the concept of  $k$ -proper index. A tree  $T$  in an edge-colored graph is a *proper tree* if any two adjacent edges of  $T$  are colored with different colors. Let  $G$  be a graph of order  $n$  and  $k$  be a fixed integer with  $2 \leq k \leq n$ . For a vertex set  $S \subseteq V(G)$ , a tree containing the vertices of  $S$  in  $G$  is called an  *$S$ -tree*. An edge-coloring of  $G$  is called a  *$k$ -proper coloring* if for every set  $S$  of  $k$  vertices in  $G$ , there exists a proper  $S$ -tree in  $G$ . The  *$k$ -proper index* of a nontrivial connected graph  $G$ , denoted by  $px_k(G)$ , is the smallest number of colors needed in a  $k$ -proper coloring of  $G$ . By definition,

$px_2(G)$  is precisely the proper connection number  $pc(G)$  for any nontrivial graph  $G$ . As a variety of nice results about  $pc(G)$  have been obtained, we in this paper only study  $px_k(G)$  for  $3 \leq k \leq n$ .

The paper is organized as follows: In Section 2, some simple observations about  $px_k(G)$  for a nontrivial graph  $G$  are stated. Meanwhile, certain necessary lemmas are also listed. In Section 3, the  $k$ -proper indices of some special graphs are determined. And for every pair of positive integers  $a, b$  with  $2 \leq a \leq b$ , a connected graph  $G$  with  $px_k(G) = a$  and  $rx_k(G) = b$  is constructed for each integer  $k$  with  $3 \leq k \leq n$ . In Section 4, the graphs with  $k$ -proper index  $n - 1$  and  $n - 2$  are characterized, respectively.

## 2. Preliminaries

We, in this section, state some observations about  $px_k(G)$  for a nontrivial graph  $G$ . Also, certain necessary lemmas are listed.

For a graph  $G$  of order  $n \geq 3$ , it follows from the definition that

$$(*) \quad pc(G) = px_2(G) \leq px_3(G) \leq px_4(G) \leq \cdots \leq px_n(G).$$

This simple property will be used frequently later.

Since any  $k$ -proper coloring of a spanning subgraph must be a  $k$ -proper coloring of its supergraph, then there exists a fundamental proposition about spanning subgraphs.

**Proposition 1.** *If  $G$  is a nontrivial connected graph of order  $n \geq 3$  and  $H$  is a connected spanning subgraph of  $G$ , then  $px_k(G) \leq px_k(H)$  for any  $k$  with  $3 \leq k \leq n$ . In particular,  $px_k(G) \leq px_k(T)$  for every spanning tree  $T$  of  $G$ .*

It has been seen in [9] that  $rx_k(G) \leq n - 1$  for any graph  $G$  of order  $n \geq 3$  and any integer  $k$  with  $3 \leq k \leq n$ . Since a rainbow tree must be a proper tree, then obviously  $px_k(G) \leq rx_k(G) \leq n - 1$ . Moreover, this simple upper bound is sharp, the graphs with  $px_k(G) = n - 1$  will be characterized in Section 4.

For any nontrivial graph  $G$ ,  $\chi'(G)$  denotes the edge-chromatic number of  $G$ . It is well-known that either  $\chi'(G) = \Delta(G)$  or  $\chi'(G) = \Delta(G) + 1$  by Vizing's Theorem, where  $\Delta(G)$ , or simply  $\Delta$ , is the maximum degree of  $G$ . Accordingly, a natural upper bound on  $px_k(G)$  with respect to these parameters follows.

**Proposition 2.** *Let  $G$  be a connected graph of order  $n \geq 3$ , the maximum degree  $\Delta(G)$  and the edge-chromatic number  $\chi'(G)$ . Then for each integer  $k$  with  $3 \leq k \leq n$ , we have*

$$px_k(G) \leq \chi'(G) \leq \Delta(G) + 1.$$

In addition, there is a classical result about the edge-chromatic number of a graph, which will be useful in the next section.

**Lemma 1** ([2]). *If  $G$  is bipartite, then  $\chi'(G) = \Delta(G)$ .*

For arbitrary  $k$  ( $k \geq 3$ ) vertices of a nontrivial graph  $G$ , any tree  $T$  containing these vertices must contain internal vertices. While for any proper tree  $T$ , there must be  $d(u)$  distinct colors assigned to the incident edges of each vertex  $u$  in  $T$ , where  $d(u)$  denotes the degree of  $u$  in  $T$ . Hence, the incident edges of any internal vertex must be assigned with at least two distinct colors to make  $T$  proper. Then the following trivial lower bound is immediate.

**Proposition 3.** *For arbitrary connected graph  $G$  of order  $n \geq 3$ , we have*

$$px_k(G) \geq 2$$

for any integer  $k$  with  $3 \leq k \leq n$ .

**Remark:** The above lower bound of  $px_k(G)$  is sharp since there exist many graphs satisfying  $px_k(G) = 2$ , as shown in Section 3. Further, we believe that it will be interesting to characterize all graphs with  $k$ -proper index 2 for specific values of  $k$ .

In any graph  $G$ , a path that contains every vertex of  $G$  is called a *Hamilton path* of  $G$ . A graph is *traceable* if it contains a Hamilton path. Bearing in mind that Inequality (\*), the following is an immediate consequence of these definitions, as well as Proposition 3.

**Proposition 4.** *Let  $G$  be a connected graph of order  $n \geq 3$ , then  $px_n(G) = 2$  if and only if  $G$  is a traceable graph.*

As mentioned before, characterizing all graphs with  $k$ -proper index 2 for specific values of  $k$  would be interesting. Proposition 4 is essentially the case of  $k = n$ . For the case of  $k = n - 1$ , there is a basic result that can be presented.

**Observation 1.** *If a connected graph  $G$  of order  $n$  satisfies  $px_{n-1}(G) = 2$ , then  $px_n(G) = 2$  if and only if  $G$  is traceable. Otherwise,  $px_n(G) = 3$ .*

It is well known that if  $G$  is a simple graph of order  $n \geq 3$  and minimum degree  $\delta \geq \frac{n-1}{2}$ , then  $G$  is traceable. Whereupon a direct corollary follows.

**Corollary 1.** *If  $G$  is a simple graph of order  $n \geq 3$  and minimum degree  $\delta \geq \frac{n-1}{2}$ , then  $px_k(G) = 2$  for each integer  $k$  with  $3 \leq k \leq n$ .*

In [9], Chartrand et al. derived the  $k$ -rainbow index of a nontrivial tree, which will be helpful in the next section.

**Lemma 2** ([9]). *Let  $T$  be a tree of order  $n \geq 3$ . For each integer  $k$  with  $3 \leq k \leq n$ ,*

$$rx_k(T) = n - 1.$$

In [3], Borozan et al. established the proper connection number of trees.

**Lemma 3** ([3]). *If  $G$  is a tree then  $pc(G) = \Delta(G)$ .*

At the end of this section, we recall several notations required in the subsequent sections.

Let  $E' \subseteq E(G)$  be a set of edges of a graph  $G$ , then  $G[E']$  is the subgraph of  $G$  induced by  $E'$ . If  $e$  is an edge of  $G$ , then  $G - e$  denotes the graph obtained from  $G$  by only deleting the edge  $e$ . If  $G$  is not complete, denote by  $G + e$  the graph obtained from  $G$  by the addition of  $e$ , where  $e$  is an edge connecting two nonadjacent vertices of  $G$ .

### 3. The $k$ -proper indices of special graphs

In this section, we determine the  $k$ -proper indices of complete graphs, cycles, wheels, trees and unicyclic graphs. Moreover, the independence of  $px_k(G)$  and  $rx_k(G)$  is given by a brief theorem.

By Proposition 4, if  $G$  is a traceable graph, then  $px_k(G) = 2$ . Obviously, the complete graphs, cycles and wheels are all traceable, thus the  $k$ -proper indices of these graphs are direct consequences of Proposition 4.

**Theorem 1.** *Let  $K_n$ ,  $C_n$  and  $W_n$  be a complete graph, a cycle and a wheel with  $n$  ( $n \geq 3$ ) vertices, respectively. Then for any integer  $k$  with  $3 \leq k \leq n$ , we have*

$$px_k(K_n) = px_k(C_n) = px_k(W_n) = 2.$$

Now we determine the  $k$ -proper index for a nontrivial tree.

**Theorem 2.** *If  $T$  is a tree of order  $n \geq 3$ , then for each integer  $k$  with  $3 \leq k \leq n$ ,*

$$px_k(T) = \Delta(T).$$

*Proof.* Firstly, since  $T$  is bipartite, then  $px_k(T) \leq \chi'(T) = \Delta(T)$  for  $3 \leq k \leq n$  by Proposition 2 and Lemma 1. On the other hand, according to Inequality (\*) and Lemma 3,  $px_k(T) \geq pc(T) = \Delta(T)$  holds naturally for  $3 \leq k \leq n$ . Therefore, we arrive at  $px_k(T) = \Delta(T)$  for any integer  $k$  with  $3 \leq k \leq n$ . ■

Combining with Proposition 1 and Theorem 2, one can check that the following proposition holds.

**Proposition 5.** *For any graph  $G$  of order  $n \geq 3$  and any integer  $k$  with  $3 \leq k \leq n$ , we have*

$$px_k(G) \leq \min\{\Delta(T) : T \text{ is a spanning tree of } G\}.$$

Since  $\Delta(T) \leq \Delta(G)$  for any spanning tree  $T$  of  $G$ , then the upper bound in Proposition 2 can be replaced by  $\Delta(G)$ .

**Proposition 6.** *Let  $G$  be a connected graph of order  $n \geq 3$  and maximum degree  $\Delta(G)$ , then*

$$px_k(G) \leq \Delta(G)$$

*for each integer  $k$  with  $3 \leq k \leq n$ .*

**Remark:** The above upper bound of  $px_k(G)$  is sharp since the equality clearly holds for any nontrivial tree.

In order to get the  $k$ -proper index of a unicyclic graph, an assistant lemma is presented.

**Lemma 4.** *Let  $G$  be a connected graph of order  $n \geq 3$  containing bridges and  $v$  be any vertex of  $G$ . Denote by  $b(v)$  the number of bridges incident with  $v$ . Set  $b = \max\{b(v) : v \in V(G)\}$ . Then for each integer  $k$  with  $3 \leq k \leq n$ , we have  $px_k(G) \geq b$ .*

*Proof.* Since for  $3 \leq k \leq n$ , it has been seen from Inequality (\*) that  $px_k(G) \geq px_3(G)$ , then we should only prove the case when  $k = 3$ . Since  $px_3(G) \geq 2$  by Proposition 3, the result is trivial when  $b = 1$  or  $2$ . Thus

we may assume that  $b \geq 3$ . Let  $u$  be a vertex with  $b(u) = b = \max\{b(v) : v \in V(G)\}$ . Let  $F = \{uw_1, uw_2, \dots, uw_b\}$  be the set of bridges incident with  $u$ . Set  $A = \{u, w_1, w_2, \dots, w_b\}$ . For any 3-set  $S = \{w_i, w_j, u\} \subseteq A$ , where  $i, j \in \{1, 2, \dots, b\}$  and  $i \neq j$ , every  $S$ -tree  $T$  must contain the edges  $uw_i$  and  $uw_j$ . Hence, the edges  $uw_i$  and  $uw_j$  receive distinct colors to make  $T$  proper, which implies that the edges  $uw_1, uw_2, \dots, uw_b$  need  $b$  distinct colors in any 3-proper coloring of  $G$ . Therefore,  $px_3(G) \geq b$ . This completes the proof.  $\blacksquare$

With the aid of Lemma 4, now we are able to deal with the  $k$ -proper index for a unicyclic graph.

**Theorem 3.** *Let  $G$  be a unicyclic graph of order  $n \geq 3$ , and maximum degree  $\Delta(G)$ . Then, for each integer  $k$  with  $3 \leq k \leq n$ ,*

$$px_k(G) = \Delta(G) - 1$$

*when  $G$  contains at most two vertices having maximum degree such that the vertices with maximum degree are all in the unique cycle of  $G$  and these two vertices (if both exist) are adjacent;*

*Otherwise,*

$$px_k(G) = \Delta(G).$$

*Proof.* Note that when  $G = C_n$ , it follows from Theorem 1 that  $px_k(G) = px_k(C_n) = 2 = \Delta(G)$  for  $3 \leq k \leq n$ . Thus in the following we assume that  $G$  is not a cycle. And assume the vertices in the unique cycle of  $G$  are  $u_1, u_2, \dots, u_g$ . Also keep in mind that  $px_k(G) \leq \Delta(G)$  for  $3 \leq k \leq n$ , which will be used later. As before, denote by  $b(v)$  the number of bridges incident with the vertex  $v$ . The discussion is divided into three cases.

**Case 1.** At first, assume that  $G$  contains a vertex, say  $u$ , satisfying

- (1) the degree of  $u$  is  $d(u) = \Delta(G)$ .
- (2)  $u$  is not in the cycle of  $G$ .

Then the incident edges of  $u$  are all bridges, i.e.,  $b(u) = d(u) = \Delta(G)$ . According to Lemma 4, we directly have  $px_k(G) \geq b(u) = \Delta(G)$  for  $3 \leq k \leq n$ . Meanwhile, Proposition 6 guarantees  $px_k(G) \leq \Delta(G)$  for  $3 \leq k \leq n$ . Accordingly, we get  $px_k(G) = \Delta(G)$  for each integer  $k$  with  $3 \leq k \leq n$  in this case.

By Case 1, if such a vertex  $u$  exists in  $G$ , then we always have  $px_k(G) = \Delta(G)$  for each integer  $k$  with  $3 \leq k \leq n$ . To avoid redundant presentation, we in the following suppose that  $G$  contains no such vertices.

**Case 2.** Now assume  $G$  simultaneously satisfies

- (3)  $G$  contains at most two vertices having maximum degree;

- (4) the vertices with maximum degree are all in the unique cycle of  $G$ ;  
(5) these two vertices (if both exist) are adjacent in  $G$ .

Then without loss of generality, suppose that  $d(u_1) = \Delta(G)$ ,  $d(u_2) \leq \Delta(G)$  and  $d(u) < \Delta(G)$  for any other vertex  $u$ . Moreover, suppose that the neighbors of  $u_1$  are  $v_1, v_2, \dots, v_{\Delta(G)-2}, v_{\Delta(G)-1} = u_2$  and  $v_{\Delta(G)} = u_g$ . Thereupon, in any 3-proper coloring  $c$  of  $G$ , based on the proof of Lemma 4, the edges  $u_1v_i$  with  $i \in \{1, 2, \dots, \Delta(G) - 2\}$  are assigned with  $\Delta(G) - 2$  distinct colors since they are all bridges incident with  $u_1$ . Without loss of generality, suppose that  $c(u_1v_1) = 1, c(u_1v_2) = 2, \dots, c(u_1v_{\Delta(G)-2}) = \Delta(G) - 2$ . Further, we claim that at least one new color is used by the edges  $u_1u_2$  and  $u_1u_g$ . For otherwise, suppose that  $c(u_1u_2) = i$  and  $c(u_1u_g) = j$  with  $i, j \in \{1, 2, \dots, \Delta(G) - 2\}$ . If  $i = j$ , then there exists no proper tree containing the vertices  $u_1, u_2$  and  $v_i$ , a contradiction. If  $i \neq j$ , then there exists no proper tree containing the vertices  $v_i, v_j$  and  $u_2$ , again a contradiction. Therefore, at least  $\Delta(G) - 2 + 1 = \Delta(G) - 1$  different colors are used by  $c$ . It follows that  $px_3(G) \geq \Delta(G) - 1$ . Thus, Inequality (\*) deduces that  $px_k(G) \geq px_3(G) \geq \Delta(G) - 1$  for each integer  $k$  with  $3 \leq k \leq n$ . On the other hand, obviously  $G - u_1u_2$  is a spanning tree of  $G$  with maximum degree  $\Delta(G) - 1$ . By Theorem 2, we know that  $px_k(G - u_1u_2) = \Delta(G - u_1u_2) = \Delta(G) - 1$  for  $3 \leq k \leq n$ . Consequently,  $px_k(G) \leq px_k(G - u_1u_2) = \Delta(G) - 1$  based on Proposition 1. To sum up, we obtain  $px_k(G) = \Delta(G) - 1$  for each integer  $k$  with  $3 \leq k \leq n$  in this case.

**Case 3.** Finally, we discuss the case when  $G$  contains at least two vertices  $u_i$  and  $u_j$  such that

- (6)  $d(u_i) = d(u_j) = \Delta(G)$ ;  
(7) both  $u_i$  and  $u_j$  are in the cycle of  $G$ ;  
(8)  $u_i$  and  $u_j$  are not adjacent in  $G$ .

Then we claim that  $px_3(G) \geq \Delta(G)$ . Assume to the contrary,  $px_3(G) \leq \Delta(G) - 1$ . Let  $c'$  be a 3-proper coloring of  $G$  using colors from  $\{1, 2, \dots, \Delta(G) - 1\}$ . Let the neighbors of  $u_i$  be  $w_1, w_2, \dots, w_{\Delta(G)-2}, w_{\Delta(G)-1} = u_{i-1}, w_{\Delta(G)} = u_{i+1}$ , and the neighbors of  $u_j$  be  $z_1, z_2, \dots, z_{\Delta(G)-2}, z_{\Delta(G)-1} = u_{j-1}, z_{\Delta(G)} = u_{j+1}$ . Similarly to Case 2, the edges  $u_iw_t$  with  $t \in \{1, 2, \dots, \Delta(G) - 2\}$  are assigned with  $\Delta(G) - 2$  distinct colors. Without loss of generality, suppose that  $c'(u_iw_1) = 1, c'(u_iw_2) = 2, \dots, c'(u_iw_{\Delta(G)-2}) = \Delta(G) - 2$ . Thus, either  $c'(u_iu_{i-1}) = c'(u_iu_{i+1}) = \Delta(G) - 1$ , or there exists at least one edge between  $u_iu_{i-1}$  and  $u_iu_{i+1}$ , say  $u_iu_{i-1}$ , such that  $c'(u_iu_{i-1}) = x_1$  with  $x_1 \in \{1, 2, \dots, \Delta(G) - 2\}$ . Similarly, the edges  $u_jz_t$  with  $t \in \{1, 2, \dots, \Delta(G) - 2\}$  also receive  $\Delta(G) - 2$  distinct colors. And for the edges  $u_ju_{j-1}$  and  $u_ju_{j+1}$ ,



either  $c'(u_j u_{j-1}) = c'(u_j u_{j+1})$ , or there exists at least one of them, say  $u_j u_{j+1}$ , such that  $c'(u_j u_{j+1}) = c'(u_j z_{x_2})$  with  $x_2 \in \{1, 2, \dots, \Delta(G) - 2\}$ .

(i) If  $c'(u_i u_{i-1}) = c'(u_i u_{i+1})$  and  $c'(u_j u_{j-1}) = c'(u_j u_{j+1})$ , then there exists no proper tree containing the vertices  $u_{i-1}$ ,  $u_{i+1}$  and  $w_1$ , a contradiction.

(ii) If  $c'(u_i u_{i-1}) = c'(u_i u_{i+1})$  and  $c'(u_j u_{j+1}) = c'(u_j z_{x_2})$  with  $x_2 \in \{1, 2, \dots, \Delta(G) - 2\}$ , then there exists no proper tree containing the vertices  $u_{j+1}$ ,  $u_j$  and  $z_{x_2}$ , a contradiction.

(iii) If  $c'(u_i u_{i-1}) = x_1$  with  $x_1 \in \{1, 2, \dots, \Delta(G) - 2\}$  and  $c'(u_j u_{j-1}) = c'(u_j u_{j+1})$ , then there exists no proper tree containing the vertices  $u_{i-1}$ ,  $u_i$  and  $w_{x_1}$ , a contradiction.

(iv) If  $c'(u_i u_{i-1}) = x_1$  with  $x_1 \in \{1, 2, \dots, \Delta(G) - 2\}$  and  $c'(u_j u_{j+1}) = c'(u_j z_{x_2})$  with  $x_2 \in \{1, 2, \dots, \Delta(G) - 2\}$ , then there exists no proper tree containing the vertices  $w_{x_1}$ ,  $u_{i-1}$  and  $z_{x_2}$ , a contradiction.

In summary, we verify that  $px_3(G) \geq \Delta(G)$ , which deduces that  $px_k(G) \geq px_3(G) \geq \Delta(G)$  for  $3 \leq k \leq n$ . Combining with  $px_k(G) \leq \Delta(G)$  for  $3 \leq k \leq n$ , we at last arrive at  $px_k(G) = \Delta(G)$  for each integer  $k$  with  $3 \leq k \leq n$  in this case.

The proof of this theorem is finished. ■

We conclude this section with a simple theorem to address the independence of  $px_k(G)$  and  $rx_k(G)$ . Notice that  $px_k(G) \leq rx_k(G)$  always holds, thus we can construct a connected graph  $G$  with  $px_k(G) = a$  and  $rx_k(G) = b$ , where  $a \leq b$ .

**Theorem 4.** *For every pair of positive integers  $a, b$  with  $2 \leq a \leq b$ , there exists a connected graph  $G$  such that  $px_k(G) = a$  and  $rx_k(G) = b$  for each integer  $k$  with  $3 \leq k \leq n$ .*

*Proof.* For each pair of positive integers  $a, b$  with  $2 \leq a \leq b$ , let  $G$  be a nontrivial tree of order  $n = b + 1$  and maximum degree  $\Delta(G) = a$ . The existence of such a tree is guaranteed by  $2 \leq a \leq b$ . Then based on Theorem 2 and Lemma 2, we know that  $px_k(G) = \Delta(G) = a$  and  $rx_k(G) = n - 1 = b$  for each integer  $k$  with  $3 \leq k \leq n$ . The proof is thus complete. ■

#### 4. Graphs with $k$ -proper index $n - 1, n - 2$

In this section, we are going to characterize the graphs whose  $k$ -proper index equals to  $n - 1$  and  $n - 2$ , respectively, where  $3 \leq k \leq n$ . First of all, we give the following lemma that will be used in the sequel.

**Lemma 5.** For  $n \geq 5$ , let  $S_n^+$  be the graph of order  $n$  obtained by adding a new edge to the  $n$ -vertices star  $S_n$ , and  $S_n^{++}$  be the graph of order  $n$  obtained by adding a new edge to  $S_n^+$ . Then we have  $px_k(S_n^{++}) \leq n-3$  for each integer  $k$  with  $3 \leq k \leq n$ .

*Proof.* Let  $V(S_n^+) = V(S_n^{++}) = \{u, v_1, v_2, \dots, v_{n-1}\}$ . Without loss of generality, set  $d_{S_n^+}(u) = d_{S_n^{++}}(u) = n-1$  and  $d_{S_n^+}(v_1) = d_{S_n^{++}}(v_2) = 2$ . Further, let  $e$  be the edge of  $S_n^{++}$  added to  $S_n^+$ . We split the remaining proof into the following two cases depending on the position of  $e$ .

**Case 1.** The edges  $e$  and  $v_1v_2$  are vertex-disjoint. Without loss of generality, let  $e = v_3v_4$ . Then,  $G' = G - uv_1 - uv_3$  is a spanning tree of  $S_n^{++}$  with maximum degree  $n-1-2 = n-3$ . It follows from Theorem 2 that  $px_k(G') = \Delta(G') = n-3$  for  $3 \leq k \leq n$ . Hence, Proposition 1 deduces that  $px_k(S_n^{++}) \leq px_k(G') = n-3$  for each integer  $k$  with  $3 \leq k \leq n$ .

**Case 2.** The edges  $e$  and  $v_1v_2$  have a common vertex. Without loss of generality, let  $e = v_2v_3$ . Then,  $G'' = G - uv_2 - uv_3$  is a spanning tree of  $S_n^{++}$  with maximum degree  $n-1-2 = n-3$ . Similarly,  $px_k(G'') = \Delta(G'') = n-3$  for  $3 \leq k \leq n$ . Hence, we can also get  $px_k(S_n^{++}) \leq px_k(G'') = n-3$  for each integer  $k$  with  $3 \leq k \leq n$ .

Combining the above two cases, now the lemma follows.  $\blacksquare$

**Theorem 5.** Let  $G$  be a connected graph of order  $n$  ( $n \geq 4$ ). Then for each integer  $k$  with  $3 \leq k \leq n$ , we have  $px_k(G) = n-1$  if and only if  $G \cong S_n$ , where  $S_n$  is the star of order  $n$ .

*Proof.* Firstly, if  $G \cong S_n$ , then by Theorem 2, we directly obtain  $px_k(G) = px_k(S_n) = \Delta(S_n) = n-1$  for  $3 \leq k \leq n$ .

Conversely, suppose  $G$  is a connected graph with  $px_k(G) = n-1$  for each integer  $k$  with  $3 \leq k \leq n$ . Since  $n-1 = px_k(G) \leq \Delta(G)$  by Proposition 6, meanwhile  $\Delta(G) \leq n-1$  holds for any simple graph of order  $n$ . Then,  $\Delta(G) = n-1$ . The hypothesis is true if  $G \cong S_n$ . If  $G \not\cong S_n$ , let  $u$  be a vertex of  $G$  with  $d(u) = \Delta(G) = n-1$ . Let  $V(G) \setminus u = \{v_1, v_2, \dots, v_{n-1}\}$  denote the set of the remaining vertices in  $G$ . Since  $G \not\cong S_n$ , there exist at least two vertices, say  $v_1$  and  $v_2$ , such that they are adjacent in  $G$ . Set  $G' = G[\bigcup_{i=1}^{n-1} uv_i] + v_1v_2$ .

Then, as  $n \geq 4$ ,  $G'$  is a unicyclic graph satisfying the conditions in Case 2 of Theorem 3 with maximum degree  $n-1$ . Hence,  $px_k(G') = \Delta(G') - 1 = n-2$  by Theorem 3. Certainly,  $G'$  is a spanning subgraph of  $G$ , therefore  $px_k(G) \leq$

$px_k(G') = n - 2$  for  $3 \leq k \leq n$  according to Proposition 1, contradicting our assumption that  $px_k(G) = n - 1$ . Consequently,  $G \cong S_n$ .

The proof is thus complete.  $\blacksquare$

**Remark:** If  $G$  is a connected graph of order  $n = 3$ , then one can check that  $px_3(G) = n - 1 = 2$  if and only if  $G \cong S_3$  or  $G \cong C_3$ .

**Theorem 6.** *Let  $G$  be a connected graph of order  $n$  ( $n \geq 5$ ). Then for each integer  $k$  with  $3 \leq k \leq n$ , we have  $px_k(G) = n - 2$  if and only if  $G \cong S_n^+$  or  $G_0$ , where  $S_n^+$  is defined in Lemma 5 and  $G_0$  is shown in Figure 1.*

*Proof.* On one hand, if  $G \cong S_n^+$ , then  $G$  is a unicyclic graph with maximum degree  $n - 1$  satisfying the conditions in Case 2 of Theorem 3. Thus  $px_k(G) = px_k(S_n^+) = \Delta(G) - 1 = n - 2$  for  $3 \leq k \leq n$ . If  $G \cong G_0$ , then  $G$  is a tree of order  $n \geq 5$  and maximum degree  $n - 2$ . Accordingly, by Theorem 2,  $px_k(G) = px_k(G_0) = \Delta(G) = n - 2$  for  $3 \leq k \leq n$ .

On the other hand, if  $px_k(G) = n - 2$ , then by Proposition 6,  $\Delta(G) \geq px_k(G) = n - 2$ , which means that  $\Delta(G) = n - 2$  or  $n - 1$ . The remaining proof is divided into two cases depending on the value of  $\Delta(G)$ .

**Case 1.**  $\Delta(G) = n - 1$ .

In this case, since  $px_k(S_n) = n - 1$  for  $3 \leq k \leq n$ , as shown before, then  $G$  must contain  $S_n^+$  as a connected spanning subgraph. If  $G \cong S_n^+$ , we have known that  $px_k(S_n^+) = n - 2$  for  $3 \leq k \leq n$ . Now suppose  $G \not\cong S_n^+$ . Then there exists a connected spanning subgraph with the form of  $S_n^{++}$  in  $G$ . Applying Proposition 1 together with Lemma 5, we arrive at  $px_k(G) \leq px_k(S_n^{++}) \leq n - 3$ , a contradiction. Hence,  $G \cong S_n^+$  in this case.

**Case 2.**  $\Delta(G) = n - 2$ .

Then  $G_0$  must be a connected spanning subgraph of  $G$ . If  $G \cong G_0$ , then  $px_k(G_0) = n - 2$  for  $3 \leq k \leq n$ . If  $G \not\cong G_0$ , then there exists at least one edge  $e \in E(G) \setminus E(G_0)$ . Thus,  $G$  contains a connected spanning subgraph isomorphic to  $G_1$ ,  $G_2$  or  $G_3$ , where  $G_1$ ,  $G_2$  and  $G_3$  are shown in Figure 1. Clearly, one can check that  $G_1$ ,  $G_2$  and  $G_3$  are all unicyclic graphs with maximum degree  $n - 2$  satisfying the conditions in Case 2 of Theorem 3. Thereupon, by Theorem 3 as well as Proposition 1, we directly get that  $px_k(G) \leq px_k(G_i) = \Delta(G_i) - 1 = n - 3$  for  $3 \leq k \leq n$  and  $i = 1, 2$  or  $3$ , which is a contradiction. Accordingly,  $G \cong G_0$  in this case.

In summary, if  $px_k(G) = n - 2$  for  $3 \leq k \leq n$ , then  $G \cong S_n^+$  or  $G \cong G_0$ . And the proof of this theorem is complete.  $\blacksquare$

**Remark:** When  $n = 4$ , except for the star  $S_4$ , other connected graphs of order 4 are all traceable. Then by Proposition 4, the  $k$ -proper indices of

these graphs equal to  $2 = n - 2$  for each integer  $k$  with  $3 \leq k \leq 4$ . While for the star  $S_4$ , we know that  $px_k(S_4) = 3$  for  $3 \leq k \leq 4$ . Consequently, we can easily claim that if  $G$  is a connected graph of order  $n = 4$ , then for each integer  $k$  with  $3 \leq k \leq 4$ ,  $px_k(G) = n - 2 = 2$  if and only if  $G \not\cong S_4$ .

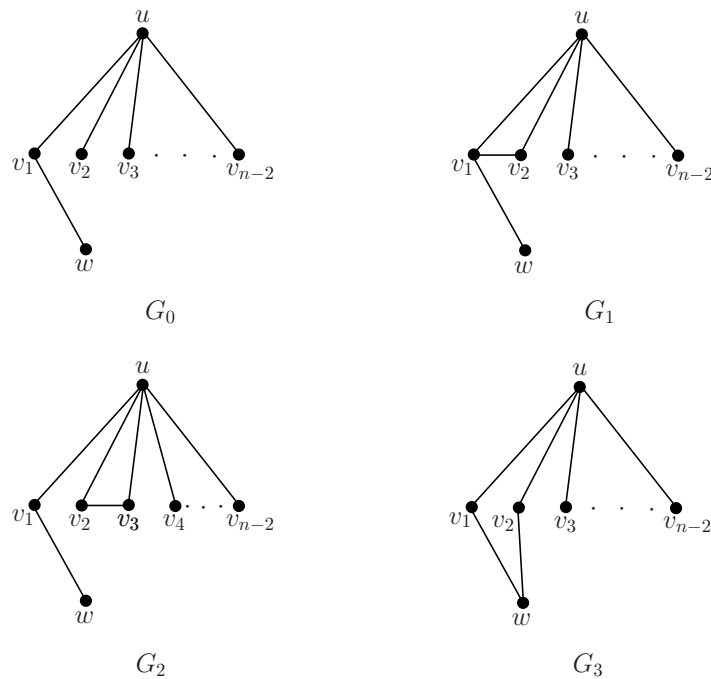


Figure 1: The graphs  $G_i$  for  $i = 0, 1, 2, 3$ .

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